

Seven (+-2) Sins of Concurrency

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In which I will show classical
concurrency problems and some
techniques of detecting and
avoiding them

I have a B.Sc. in CS and Statistics,
OCP, 10 years of production IT
experience and I'm an Oracle
Ace. So I know what I'm talking
about.

But you don't have to trust me – I
have scripts that prove everything I
say.



Since 1967
computers
can walk and
chew gum at
the same
time

Programs need to learn to share

Example: Shared Bank Account

```
create or replace procedure update_account (p_id
    number,p_amount number) as
    n number;
begin
    SELECT amount into n FROM bank_account WHERE
    id=p_id;
    UPDATE bank_account SET amount = n+p_amount
    WHERE id=p_id;
end;
```

```
SQL> exec deposit(1,500)
```

```
SQL> commit;
```

```
SQL> exec withdraw(1,-500)
```

```
SQL> commit;
```

```
SQL> select amount from bank_account;
```

```
AMOUNT
```

```
-----
```

```
-500
```


Sin #1

Race Condition



Can your code share?

Are you 100% sure?

Does this look familiar?

```
spool XXX_drop_db_links.sql
```

```
select 'drop database link '||OBJECT_NAME||';'  
from obj  
where OBJECT_TYPE='DATABASE LINK';
```

```
spool off
```

```
@XXX_drop_db_links.sql
```

Sin #2

Ostrich Algorithm



Few words about critical sections
and mutual exclusion

Laws of Good Concurrency

- No two processes will be in their critical section at same time
- No assumptions about number or speed of CPUs
- No process outside the critical section may block other processes
- No process will wait forever to enter critical section

Mutual Exclusion in Oracle

Locks and latches and mutexes,
oh my!

User Defined Locks

```
dbms_lock.allocate_unique (  
    lockname      => 'Synchronize',  
    lockhandle    => m_handle  
);  
  
n1 := dbms_lock.request (  
    lockhandle      => m_handle,  
    lockmode        => dbms_lock.x_mode,  
    timeout         => dbms_lock.maxwait,  
    release_on_commit => true  
);
```



```
dbms_lock.allocate_unique('Synchronize',m_handle);
dbms_lock.request(m_handle,dbms_lock.x_mode,
    dbms_lock.maxwait,false);
spool XXX_drop_db_links.sql

select 'drop database link '||OBJECT_NAME||';'
from obj where OBJECT_TYPE='DATABASE LINK';

spool off

@XXX_drop_db_links.sql

dbms_lock.release(m_handle);
```

Another Race

```
select max(id) into max_id from my_table;
```

```
insert into my_table values (max_id+1,some_data);
```

```
commit;
```



Protecting the critical section - I

```
select max(id) into max_id from my_table for update;  
  
insert into my_table values (max_id+1, some_data);  
  
commit;
```

ERROR at line 1:

ORA-01786: FOR UPDATE of this query expression is
not allowed



Protecting the critical section - II

```
select id into max_id from my_table where id=(select  
    max(id) from my_table) for update;
```

```
insert into my_table values (max_id+1,some_data);
```

```
commit;
```



Protecting the critical section - III

```
select max_id into p_max_id from extra_table for  
update;
```

```
insert into my_table values (max_id+1, some_data);
```

```
update extra_table set max_id=max_id+1;
```

```
commit;
```

Sin #3

Solving the race condition led to
serialization

The right way to do it:

```
create sequence my_table_seq start with 1  
    increment by 1 cache 20;
```

```
insert into my_table  
    (my_table_seq.nextval, some_data);  
commit;
```



Quick Review

"Insanity: Doing the same thing over and over again and expecting different results."

Albert Einstein.

Classical Concurrency Problems



Dining Philosophers



```
-- number of philosophers
select count(*) into N from sticks;
think();

update sticks set owner=philosopher_id where
    s_id=p_id; -- take right fork
update sticks set owner=philosopher_id where
    s_id=mod(p_id+1,N); -- take left fork

eat(); -- nom nom nom
commit; -- put down forks
```

ORA-00060: Deadlock detected



Sin #4

```
think ();
```

```
update sticks set owner=in_p_id where s_id=in_p_id;
```

```
select s_id into r_s from sticks where  
s_id=mod(in_p_id+1,N) for update nowait;
```

```
update sticks set owner=in_p_id where  
s_id=mod(in_p_id+1,N);
```

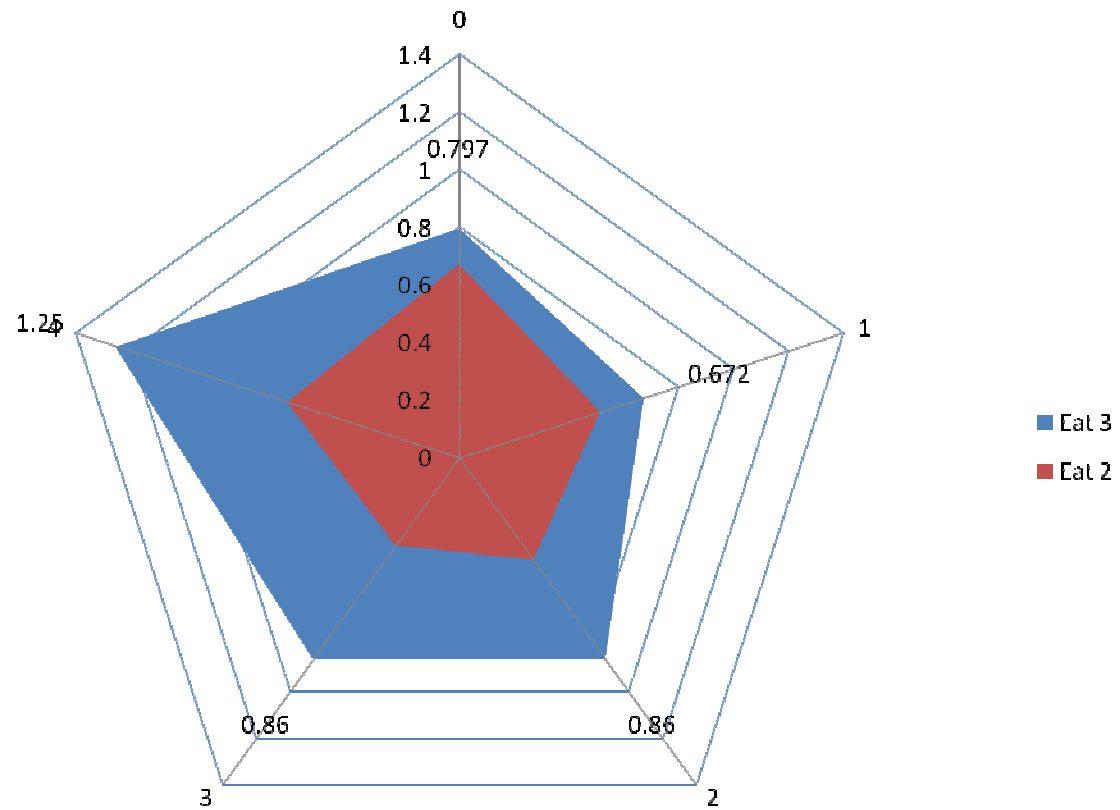
```
eat ();
```

```
commit;
```

```
exception
```

```
when resource_busy then  
rollback;
```

Starvation



Sin #5

```
think();
```

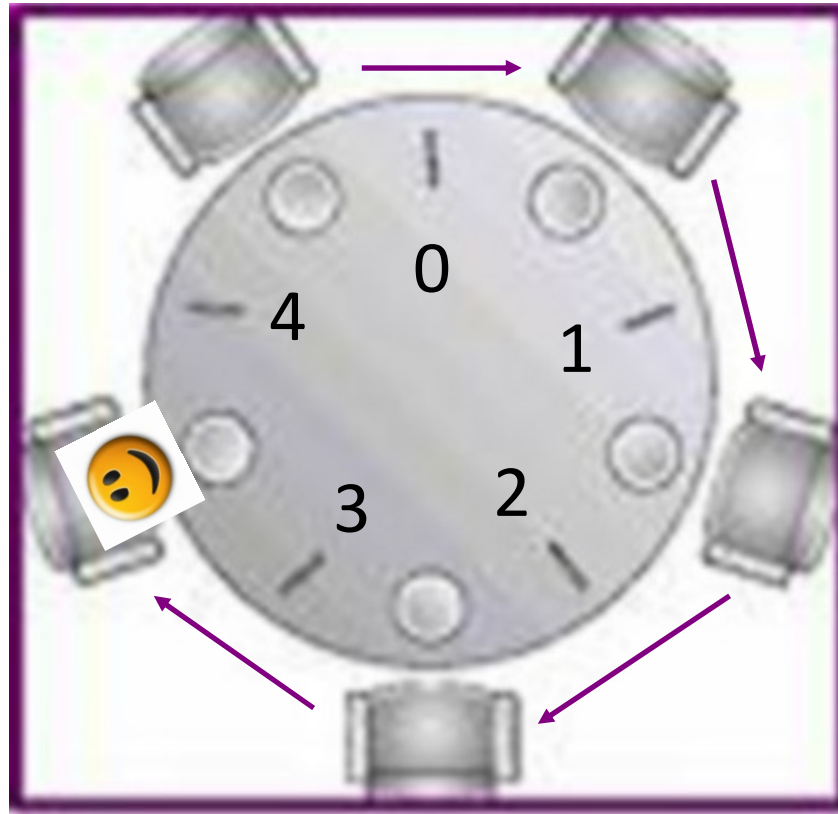
```
update sticks set owner=in_p_id where  
  s_id=least(in_p_id,mod(in_p_id+1,N));
```

```
update sticks set owner=in_p_id where  
  s_id=greatest(in_p_id,mod(in_p_id+1,N));
```

```
eat();
```

```
commit;
```


Partial Hierarchy Solution



Or just index your foreign keys!



Quick Review

Barbershop Queue



Generating customers

```
update customers set
needs_cut=1,entered_shop=systimestamp
where id in (
    select id from
        (select id from customers
         where needs_cut=0
         order by dbms_random.random)
    where
        rownum<=(dbms_random.value*(p_avg_customers_per_sec*2+1)));

commit;
dbms_lock.sleep(1);
```

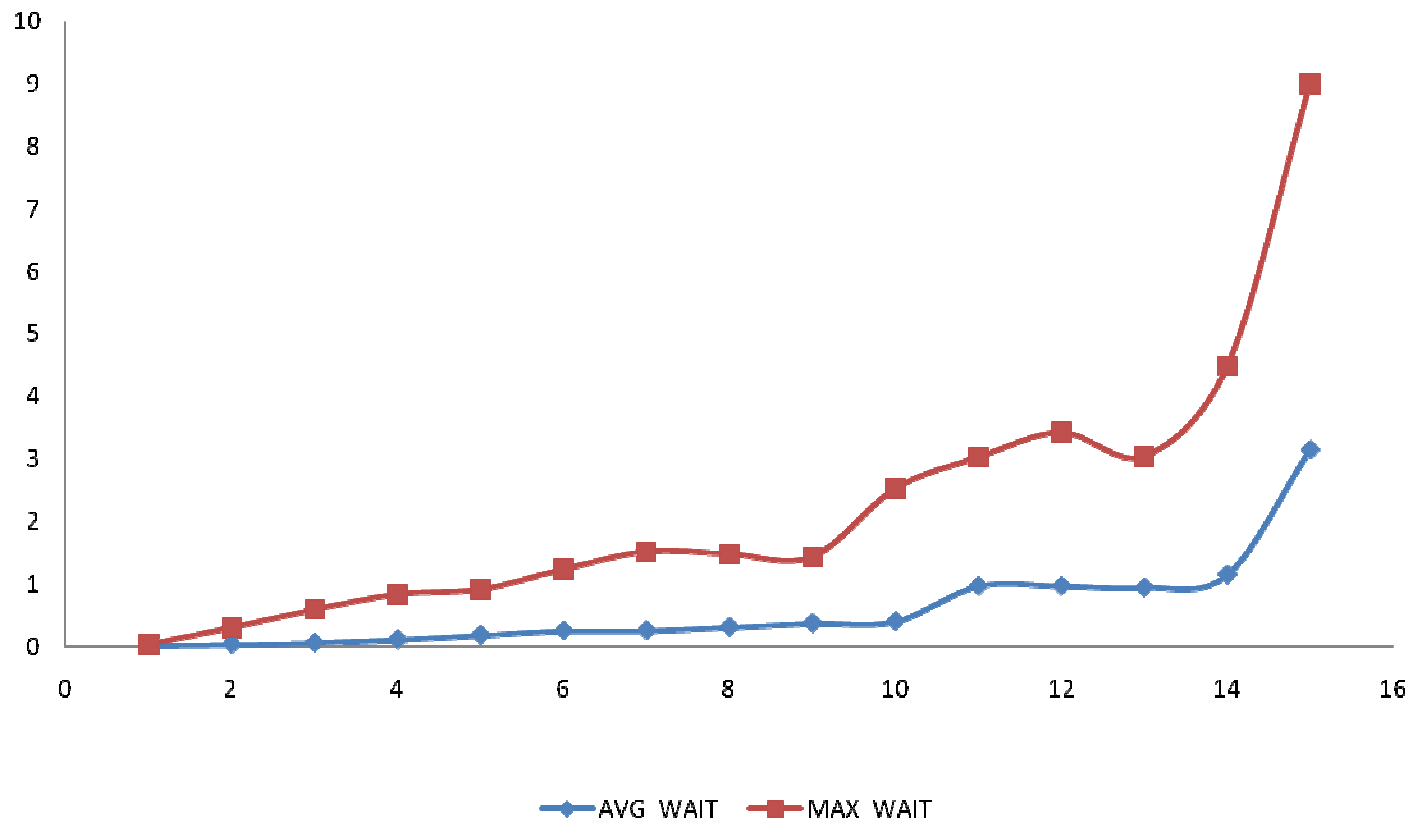
Each barber does:

```
cursor c is select * from customers where  
  needs_cut=1 order by entered_shop for update skip  
  locked;
```

...

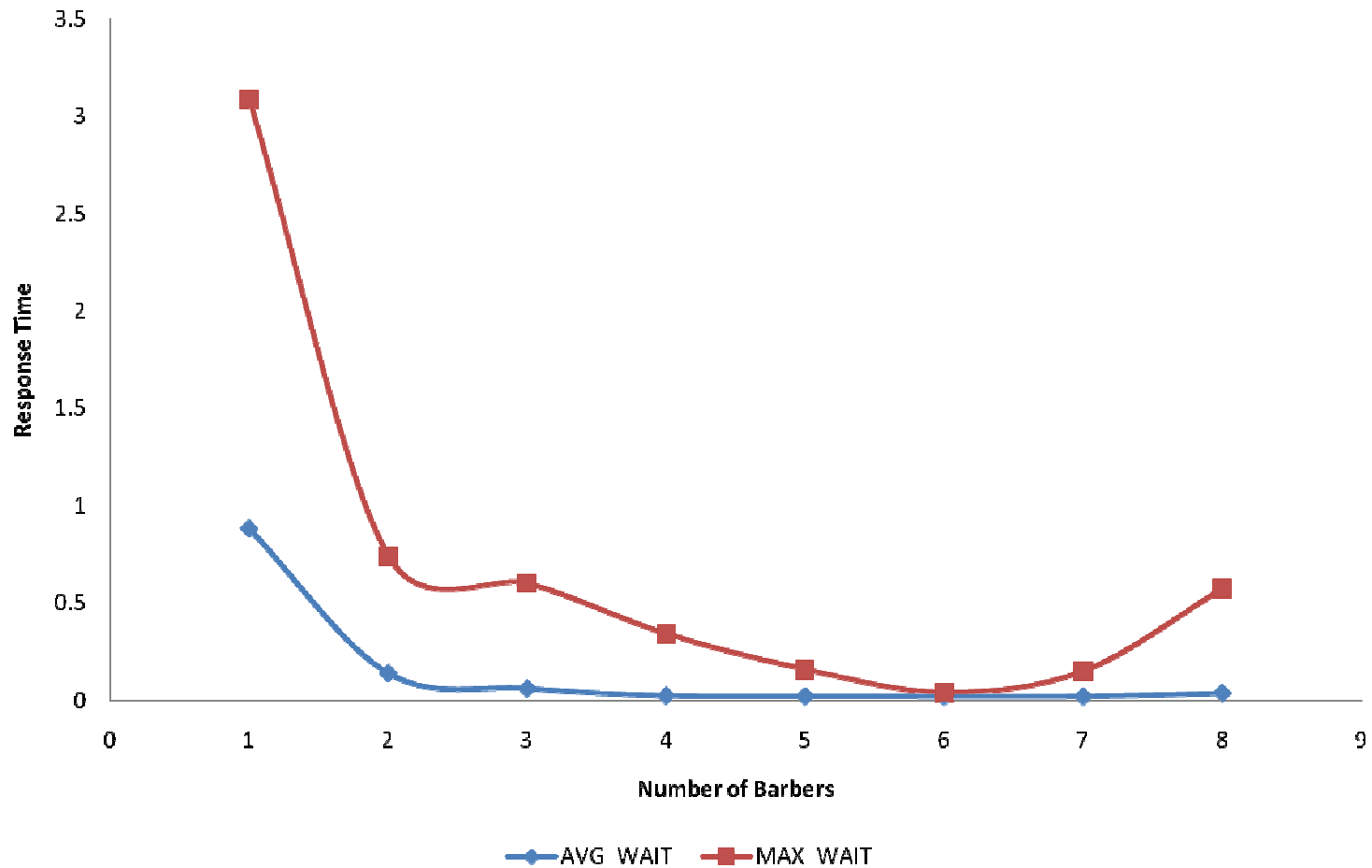
```
fetch c into l_rec;  
exit when c%NOTFOUND;  
cut_hair(dbms_random.value*p_avg_cut_time*2);  
finish_work(l_rec.id);
```

3 Barbers, Haircut in 0.3 seconds



Sin #6

About 3 concurrent customers



Sin #7

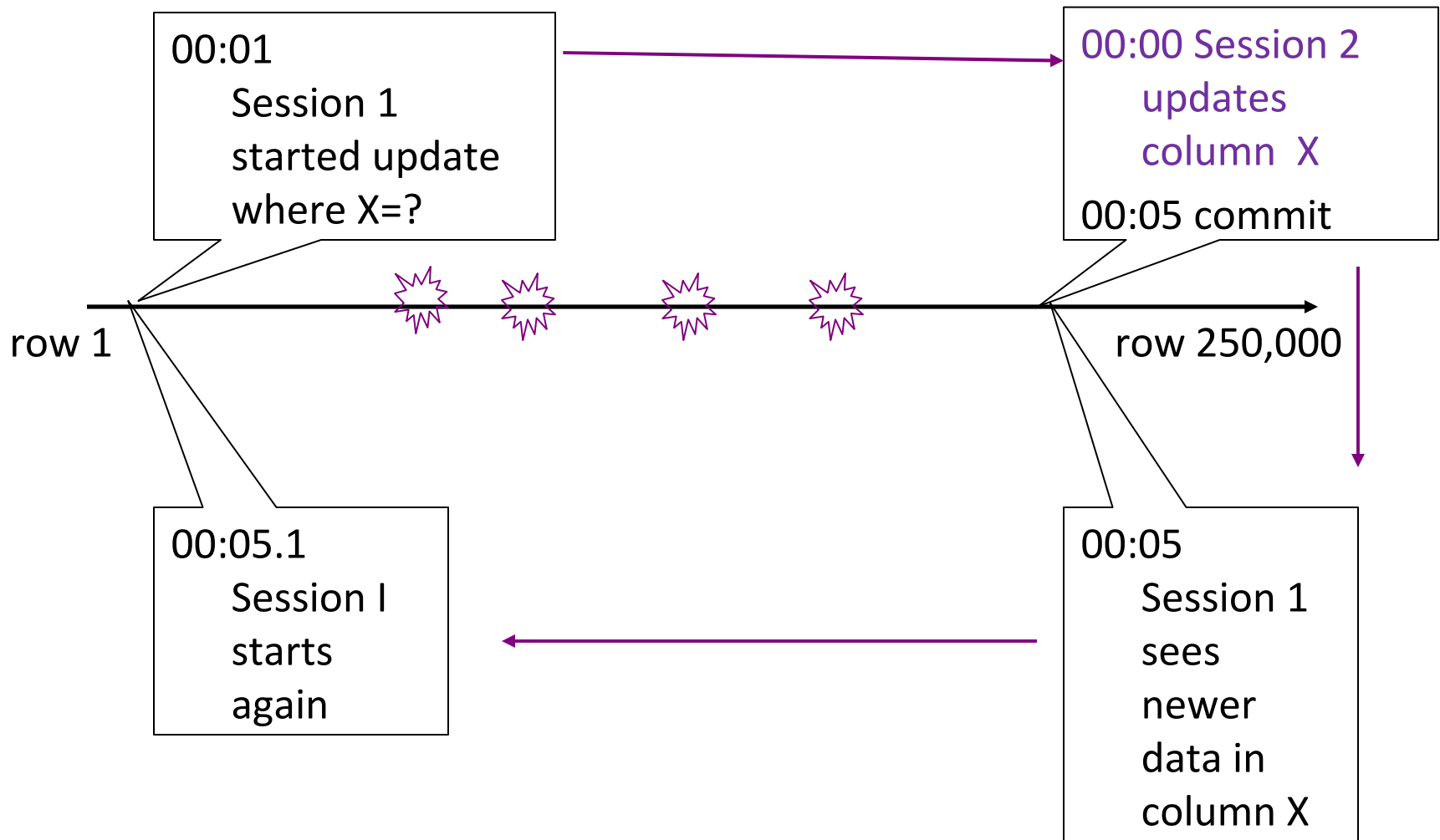


Quick Review

Oracle Concurrency Problems

Because consistency has a price

Non-transactional changes



Forgetting the extra IO

